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Enabling Fast Forwarding in Hybrid Software-Defined Networks

Yijun Sun, Zehua Guo, Songshi Dou, Junjie Zhang, Changlin Li, and Xiang Ouyang

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Outline

Introduction

- Problem analysis and solution
- **3** Preliminary result





SDN and Hybrid SDN





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Software-Defined Networking (SDN)

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Hybrid SDN

1 SDN and Hybrid SDN





Hybrid-SDN

Key issue:

maximize network programmability for SDN by upgrading a given number of legacy devices to SDN devices.

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1 Layer-2 network

- Layer-2 network:
- MAC address
- VLAN and STP
- Campus and enterprise network





1 **Opportunity**





Network topology

Two virtual networks

Flow f: 1->5 in virtual network 2

Upgrade benefit: reducing the forwarding path's length

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Problem analysis and solution

2 **Problem analysis-Formulation**

L: the number of virtual networks, each of which is a spanning tree
N: the number of nodes at each virtual network
M: the number of nodes to be upgraded
dis(i,j): the shortest path's length between node i and node j among different virtual networks

Upgrade = add edge across the same node at different virtual networks

Challenge1:

The sequence of node upgrade matters.

(Greedy selection may not be the optimal solution)

Challenge2:

The solution space is huge. (**N** could be thousands, and **L** could be dozens)

2 Reducing the path length

Upgrade the first node k

 $dis_{new}(i,j) = \min\{\min_{u=1..L} \{dis_{old}^{u}(i,k)\} + \min_{u=1..L} \{dis_{old}^{u}(k,j)\}, \min_{u=1..L} \{dis_{old}^{u}(i,j)\}\}$ Upgrade another node p

 $dis_{new}(i,j) = \min\{dis_{old}(i,k) + dis_{old}(k,j), dis_{old}(i,j)\}$

Upgrade benefit = $\sum_{i=1}^{N} \sum_{j=1}^{N} (dis_{old}(i,j)) - \sum_{i=1}^{N} \sum_{j=1}^{N} (dis_{new}(i,j))$

(1)Start (2)Calculate the upgrade benefit of nodes from S1 (3)Select the node with the maximum upgrade benefit from S1 (4)Upgrade node j, which is selected at step 3 (5)Delete node j in S1 and add node j into S2 $(6) dis_{old} = dis_{new}$ Floyd or Dijkstra is too (7)If |S2| < M, go to step 1 slow since it calculates too many meaningless (8)End repeated paths.

S1 records nodes that have not been upgraded S2 records all upgraded nodes

Shortcutter's complexity:

 $O(n^2)$

Degree-centric: it sequentially upgrades nodes following the decreasing order of these nodes' degrees

Betweenness-centric: it sequentially upgrades nodes following the decreasing order of these nodes' betweenness centrality.

Closeness-centric: it sequentially upgrades nodes following the decreasing order of these nodes 'closeness centrality.

Simulation setup

Settings	Value
Power law degree distribution-Number of nodes	30
Power law degree distribution-Number of edges connected to the new node	1
Power law degree distribution-Probability of adding a ring	0
Number of virtual networks	5

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Result

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